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# **Extensive Form Games**

# **Definition of Extensive Form Game**

a finite game tree X with nodes  $x \in X$ 

nodes are partially ordered and have a single root (minimal element)

terminal nodes are  $z \in Z$  (maximal elements)



Players and Information Sets

player 0 is nature

information sets  $h \in H$  are a partition of  $X \setminus Z$ 

each node in an information set must have exactly the same number of immediate followers

each information set is associated with a unique player who "has the move" at that information set

 $H_i \subset H$  information sets where *i* has the move

### More Extensive Form Notation

information sets belonging to nature  $h \in H_0$  are singletons

A(h) feasible actions at  $h \in H$ 

each action and node  $a \in A(h), x \in h$  is associated with a unique node that immediately follows x on the tree

each terminal node has a payoff  $r_i(z)$  for each player

by convention we designate terminal nodes in the diagram by their payoffs

Example: a simple simultaneous move game



### **Behavior Strategies**

a *pure strategy* is a map from information sets to feasible actions  $s_i(h_i) \in A(h_i)$ 

a *behavior strategy* is a map from information sets to probability distributions over feasible actions  $\pi_i(h_i) \in P(A(h_i))$ 

*Nature's move* is a behavior strategy for Nature and is a fixed part of the description of the game

We may now define  $u_i(\pi)$ 

*normal form* are the payoffs  $u_i(s)$  derived from the game tree



U

D

L	R
1,1	2,2
3,3	4,4

7

Kuhn's Theorem:

every mixed strategy gives rise to a unique behavior strategy

The converse is NOT true



```
1 plays .5 U
behavior: 2 plays .5L at U; .5L at R
mixed: 2 plays .5(LL),.5(RR)
2 plays .25(LL),.25(RL),.25(LR),.25(RR)
```

however: if two mixed strategies give rise to the same behavior strategy, they are *equivalent*, that is they yield the same payoff vector for each opponents profile  $u(\sigma_i, s_{-i}) = u(\sigma_i^{,i}, s_{-i})$ 

## **Subgame Perfection**

some games seem to have too many Nash equilibria

### **Ultimatum Bargaining**

player 1 proposes how to divide \$10 in pennies

player 2 may accept or reject

Nash: any proposal by player 1 with all poorer proposals rejected and equal or better proposals accepted

# Chain Store



	out	in
fight	2*,0*	-1,-1
give in	2,0	1*,1*

Subgame Perfection

A subgame perfect Nash Equilibrium is a Nash equilibrium in every subgame

A subgame starts at a singleton information set



equilibria:

UR is subgame perfect

D and .5 or more L is Nash but not subgame perfect

### **Application to Rubinstein Bargaining**

the pie division game: there is one unit of pie; player 1 demands  $p_1$ 

player 2 accepts or rejects

if player 2 rejects one period elapses, then the roles are reversed, with player 2 demanding  $p_2$ 

common discount factor  $0 < \delta < 1$ 

**Nash equilibrium**: player 1 gets all pie, rejects all positive demands by player 2; player 2 indifferent and demands nothing

conversely: player 2 gets all the pie

wait 13 periods then split the pie 50-50; if anyone makes a positive offer during this waiting period, reject then revert to the equilibrium where the waiting player gets all the pie

**subgame perfection:** one player getting all pie is not an equilibrium: if your opponent must wait a period to collect all pie, he will necessarily accept demand of  $1 - \delta - \varepsilon$  today, since this give him  $\delta + \varepsilon$  in present value, rather than  $\delta$  the present value of waiting a period

Rubinstein's Theorem:

there is a unique subgame perfect equilibrium

players always make the same demands, and if they demand no more than the equilibrium level their demands are accepted

to compute the unique equilibrium observe that a player may reject an offer, wait a period, make the equilibrium demand of *p* and have it accepted, thus getting  $\delta p$  today; this means the opposing player may demand up to  $1 - \delta p$  and have the demand accepted; the equilibrium condition is

$$p = 1 - \delta p$$
 or  $p = \frac{1}{1 + \delta}$ 

notice that the player moving second gets

 $\frac{\delta}{1+\delta}$  and that as  $\delta \rightarrow 1$  the equilibrium converges to a 50-50 split

a problem: if offers are in pennies, subgame perfect equilibrium is not unique

How to prove the equilibrium is unique:

let p be such that any higher demand will be rejected in every equilibrium

let q be such that any lower demand will be accepted in every equilibrium

if you accept p you get 1 - p versus at least  $\delta q$  by rejecting, so  $1 - \delta q$ or less will be rejected in any equilibrium and  $p \le 1 - \delta q$ 

if you accept q you get 1 - q versus at most  $\delta p$  by rejecting so  $(1 - q)/\delta$  will be accepted in any equilibrium and  $p \ge (1 - q)/\delta$ 

moreover  $p \ge q$ 



# **Trembling Hand Perfection**

Selten Game



	L	R
U	-1,-1	2,0
D	1,1	1,1

subgame perfect

equilibria:

UR is subgame perfect

D and .5 or more L is Nash but not subgame perfect

can also solve by weak dominance or by trembling hand perfection

### Example of Trembling Hand not Subgame Perfect





Here Ld,D is trembling hand perfect but not subgame perfect

definition of the agent normal form

each information set is treated as a different player, e.g. 1a, 1b if player 1 has two information sets; players 1a and 1b have the same payoffs as player 1

extensive form trembling hand perfection is trembling hand perfection in the agent normal form

what is sequentiality??

# **Robustness – The Selten Game**



### genericity in normal form

	L	R
U	-1,-1	2**,0**
D	<b>1**,1*(±</b> ε)	1,1

# Self Confirming Equilibrium

- $s_i \in S_i$  pure strategies for *i*;  $\sigma_i \in \Sigma_i$  mixed
- $H_i$  information sets for *i*

 $\overline{\mathit{H}}(\sigma)$  reached with positive probability under  $\sigma$ 

 $\pi_i \in \Pi_i$  behavior strategies

 $\hat{\pi}(h_i | \sigma_i)$  map from mixed to behavior strategies

 $\hat{\rho}(\pi)$ ,  $\hat{\rho}(\sigma) \equiv \hat{\rho}(\hat{\pi}(\sigma))$  distribution over terminal nodes

 $\mu_i$  a probability measure on  $\Pi_{-i}$ 

 $u_i(s_i|\mu_i)$  preferences

 $\Pi_{-i}(\sigma_{-i}|J) \equiv \{\pi_{-i}|\pi_i(h_i) = \hat{\pi}(h_i|\sigma_i), \forall h_i \in H_{-i} \cap J\}$ 

Notions of Equilibrium

Nash equilibrium

a mixed profile  $\sigma$  such that for each  $s_i \in \text{supp}(\sigma_i)$  there exist beliefs  $\mu_i$  such that

- $s_i$  maximizes  $u_i(\cdot | \mu_i)$
- $\mu_i(\Pi_{-i}(\sigma_{-i}|H)) = 1$

Unitary Self-Confirming Equilibrium

•  $\mu_i(\Pi_{-i}(\sigma_{-i}|\overline{H}(\sigma))) = 1$ 

(=Nash with two players)





 $A_1, A_2$  is self-confirming, but not Nash

any strategy for 3 makes it optimal for either 1 or 2 to play down but in self-confirming, 1 can believe 3 plays R; 2 that he plays L Heterogeneous Self-Confirming equilibrium

• 
$$\mu_i(\Pi_{-i}(\sigma_{-i}|\overline{H}(s_i,\sigma))) = 1$$

Can summarize by means of "observation function"

 $J(s_i, \sigma) = H, \overline{H}(\sigma), \overline{H}(s_i, \sigma)$ 

#### Public Randomization



Remark: In games with perfect information, the set of heterogeneous self-confirming equilibrium payoffs (and the probability distributions over outcomes) are convex

## **Ultimatum Bargaining Results**



### Raw US Data for Ultimatum

X	Offers	Rejection Probability
\$2.00	1	100%
\$3.25	2	50%
\$4.00	7	14%
\$4.25	1	0%
\$4.50	2	100%
\$4.75	1	0%
\$5.00	13	0%
	27	
	<u> </u>	

US \$10.00 stake games, round 10

Trials	Rnd	Cntry	Case	Expect	ed Loss	Max	Ratio	
		Stake		PI 1	PI 2	Both	Gain	
27	10	US	Н	\$0.00	\$0.67	\$0.34	\$10.00	3.4%
27	10	US	U	\$1.30	\$0.67	\$0.99	\$10.00	9.9%
10	10	USx3	Н	\$0.00	\$1.28	\$0.64	\$30.00	2.1%
10	10	USx3	U	\$6.45	\$1.28	\$3.86	\$30.00	12.9%
30	10	Yugo	Н	\$0.00	\$0.99	\$0.50	\$10?	5.0%
30	10	Yugo	U	\$1.57	\$0.99	\$1.28	\$10?	12.8%
29	10	Jpn	Н	\$0.00	\$0.53	\$0.27	\$10?	2.7%
29	10	Jpn	U	\$1.85	\$0.53	\$1.19	\$10?	11.9%
30	10	Isrl	Н	\$0.00	\$0.38	\$0.19	\$10?	1.9%
30	10	Isrl	U	\$3.16	\$0.38	\$1.77	\$10?	17.7%
	WC		Н			\$5.00	\$10.00	50.0%

Rnds=Rounds, WC=Worst Case, H=Heterogeneous, U=Unitary

#### Comments on Ultimatum

- every offer by player 1 is a best response to beliefs that all other offers will be rejected so player 1's heterogeneous losses are always zero.
- big player 1 losses in the unitary case
- player 2 losses all knowing losses from rejected offers; magnitudes indicate that subgame perfection does quite badly
- as in centipede, tripling the stakes increases the size of losses a bit less than proportionally (losses roughly double).

### Centipede Game: Palfrey and McKelvey



Numbers in square brackets correspond to the observed conditional probabilities of play corresponding to rounds 6-10, stakes 1x below.

This game has a unique self-confirming equilibrium; in it player 1 with probability 1 plays  $T_1$ 

Summary of Experimental Results

Trials /	Rnds	Stake	Ca se	Expected Loss			Max	Ratio
Rnd				PI 1	PI 2	Both	Gain	
29*	6-10	1x	Η	\$0.00	\$0.03	\$0.02	\$4.0 0	0.4%
29*	6-10	1x	U	\$0.26	\$0.17	\$0.22	\$4.0 0	5.4%
	WC	1x	Η			\$0.80	\$4.0 0	20.0 %
29	1-10	1x	Η	\$0.00	\$0.08	\$0.04	\$4.0 0	1.0%
10	1-10	4x	Η	\$0.00	\$0.28	\$0.14	\$16. 00	0.9%

### Rnds=Rounds, WC=Worst Case, H=Heterogeneous, U=Unitary \*The data on which from which this case is computed is reported above.

### Comments on Experimental Results

- heterogeneous loss per player is small; because payoffs are doubling in each stage, equilibrium is very sensitive to a small number of player 2's giving money away at the end of the game.
- unknowing losses far greater than knowing losses
- quadrupling the stakes very nearly causes  $\overline{\epsilon}$  to quadruple
- theory has substantial predictive power: see WC
- losses conditional on reaching the final stage are quite large-inconsistent with subgame perfection. McKelvey and Palfrey
  estimated an incomplete information model where some "types" of
  player 2 liked to pass in the final stage. This cannot explain many
  players dropping out early so their estimated model fits poorly.