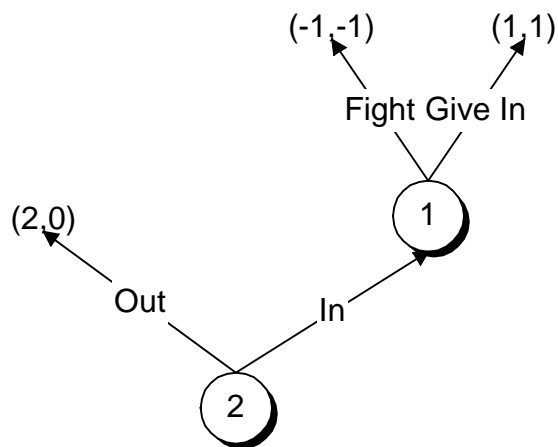


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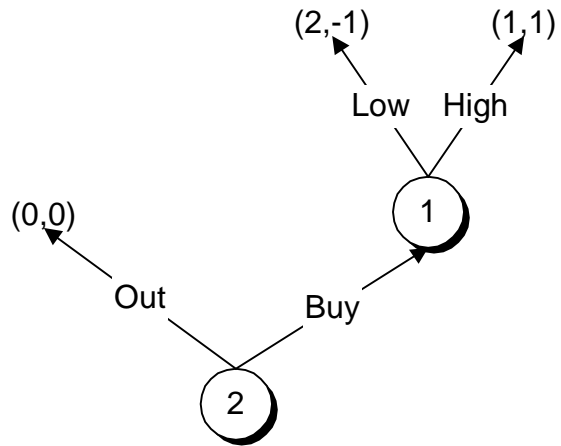
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# Reputation

## *Extensive Form Examples*



Chain Store Game



## Quality Game

## *Simultaneous Move Examples*

### *Modified Chain Store*

	out	in
fight	$2-\varepsilon, 0$	$-1,-1$
give in	$2,0$	$1,1^{**}$

## *Inflation Game*

	Low	High
Low	0,0	-2,-1
High	1,-1	-1,0

Inflation Game: LR=government, SR=consumers  
consumer preferences are whether or not they guess right

	Low	High
Low	0,0	0,-1
High	-1,-1	-1,0

with a hard-nosed government

## ***The Model***

multiple types of long-run player  $\omega \in \Omega$

$\Omega$  is a countable set of types

type is fixed forever (does not change from period to period)

$u^1(a, \omega)$  utility depends on type

strategy  $\sigma^1(h, \omega)$  depends on type

types are privately known to long-run player, not known to short run player

strategy  $\sigma^2(h)$  does not depend on type

$\mu$  probability distribution over  $\Omega$  commonly known short-run player prior over types

## *Truly Committed Types*

type  $\omega(a^1)$  has a dominant strategy to play  $a^1$  in the repeated game:

$$u^1(\tilde{a}^1, a^2, \omega(a^1)) = \begin{cases} 1 & \tilde{a}^1 = a^1 \\ 0 & \tilde{a}^1 \neq a^1 \end{cases}$$

for example

A truly committed type “can't be bargained with...can't be reasoned with...doesn't feel pity, or remorse, or fear.”

Let  $n(\omega)$  be the least utility received by a type  $\omega$  in any Nash equilibrium

let  $a^1*$  be a pure strategy Stackelberg strategy for type  $\omega_0$ , with corresponding utility

$$u^1* = \max_{\alpha^1} \min_{\alpha^2 \in BR(\alpha^1)} u^1(\alpha^1, \alpha^2, \omega_0)$$

*Theorem:* Fix  $\omega_0$  with  $\mu(\omega_0) > 0$ , and. Let  $\omega^* \equiv \omega(a^1*)$ , and suppose that  $\mu^* \equiv \mu(\omega^*) > 0$ . Then there is a constant  $k(\mu^*)$  otherwise independent of  $\mu, \Omega$  such that

$$n(\omega) \geq \delta^{k(\mu^*)} u^1* + (1 - \delta^{k(\mu^*)}) \underline{u}^1$$



## *Proof*

define  $\pi_t^*$  to be the probability at the beginning of period  $t$  by the short-run player that he is facing type  $\omega^*$

Let  $N(\pi_t^* \leq \bar{\pi})$  be the number of times  $\pi_t^* \leq \bar{\pi}$

Lemma 1: Suppose that LR plays  $a^1^*$  always. Then for any history  $h$  that has positive probability

$$pr(N(\pi_t^* \leq \bar{\pi}) > \log \mu^* / \log \bar{\pi} | h) = 0$$

Lemma 2: There is  $\bar{\pi} < 1$  such that if  $\pi_t^* > \bar{\pi}$  the SR player plays a best response to  $a^1^*$

- Why do these Lemma's imply the theorem?
- Why is Lemma 2 true?

## *Proof of Lemma 1*

Bayes Law

$$\pi(\omega^*|h_t) = \frac{\pi(\omega^*|h_{t-1})\pi(h_t|\omega^*,h_{t-1})}{\pi(h_t|h_{t-1})}$$

given  $h_{t-1}$  player 1 and 2 play independently

$$\pi(\omega^*|h_t) = \frac{\pi(\omega^*|h_{t-1})\pi(h_t|\omega^*,h_{t-1})}{\pi(h_t^1|h_{t-1})\pi(h_t^2|h_{t-1})}$$

since player 1's type isn't known to player 2

$$\pi(\omega^*|h_t) = \frac{\pi(\omega^*|h_{t-1})\pi(h_t|\omega^*,h_{t-1})}{\pi(h_t^1|h_{t-1})\pi(h_t^2|\omega^*,h_{t-1})}$$

since player 1's strategy is to always play  $a^1$  \*  $\pi(h_t^1|\omega^*,h_{t-1}) = 1$  so

$$\begin{aligned}\pi(\omega^*|h_t) &= \frac{\pi(\omega^*|h_{t-1})\pi(h_t^2|\omega^*,h_{t-1})}{\pi(h_t^1|h_{t-1})\pi(h_t^2|\omega^*,h_{t-1})} \\ &= \frac{\pi(\omega^*|h_{t-1})}{\pi(h_t^1|h_{t-1})} = \frac{\pi(\omega^*|h_{t-1})}{\pi_t^*}\end{aligned}$$

the conclusion:

$$\pi(\omega^*|h_t) = \frac{\pi(\omega^*|h_{t-1})}{\pi_t^*}$$

- what does this say?

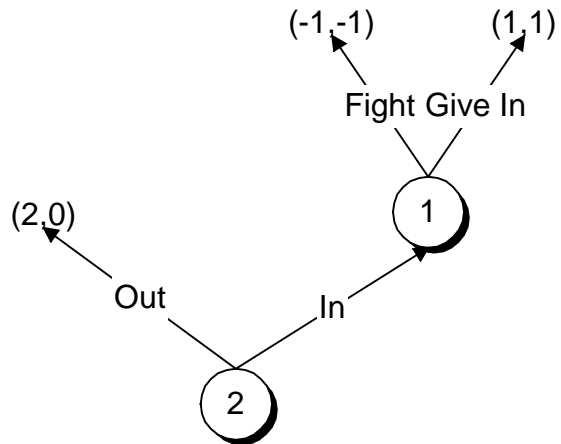
the Lemma now derives from the fact that  $\pi(\omega^*|h_t) \leq 1$

## ***Observational Equivalence***

$\rho(y|\alpha)$  outcome function

$\alpha^2 \in W(\alpha^1)$  if there exists  $\tilde{\alpha}^1$  such that  $\rho(\cdot|\tilde{\alpha}^1, \alpha^2) = \rho(\cdot|\alpha^1, \alpha^2)$  and  $\alpha^2 \in BR(\tilde{\alpha}^1)$

$$u^{1*} = \max_{\alpha^1} \min_{\alpha^2 \in W(\alpha^1)} u^1(\alpha^1, \alpha^2, \omega_0)$$



Chain Store Game

strategies that are observationally equivalent

	out	in	mixed
fight	all	fight	fight
give	all	give	give
mixed	all	mixed	mixed

weak best responses

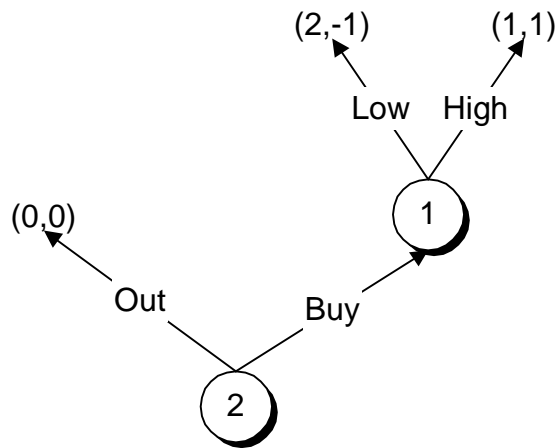
fight: out

give: in, out

mixed: in, out?

Best case fight:out so  $u^{1*} = 2$





## Quality Game

strategies that are observationally equivalent

	out	buy	mixed
hi	all	hi	hi
lo	all	lo	lo
mixed	all	mixed	mixed

weak best responses

hi: in, out

lo: out

mixed: in?, out

in every case out is a weak best response so  $u^{1*} = 0$

## ***Moral Hazard and Mixed Commitments***

$\rho(y|\alpha)$  outcome function

expand space of types to include types committed to mixed strategies:  
leads to technical complications because it requires a continuum of types

$p(h_{t-1})$  probability distribution over outcomes conditional on the history  
(a vector)

$p^+(h_{t-1})$  probability distribution over outcomes conditional on the history  
and the type being in  $\Omega^+$

**Theorem:** for every  $\varepsilon > 0, \Delta_0 > 0$  and set of types  $\Omega^+$  with  $\mu(\Omega^+) > 0$  there is a  $K$  such that if  $\Omega^+$  is true there is probability less than  $\varepsilon$  that there are more than  $K$  periods with

$$\|p^+(h_{t-1}) - p(h_{t-1})\| > \Delta_0$$

look for tight bounds

let  $\underline{n}, \bar{n}$  be best and worst Nash payoffs to LR

try to get

$$\liminf_{\delta \rightarrow 1} \underline{n}(\omega) = \limsup_{\delta \rightarrow 1} \bar{n}(\omega) = \max_{\alpha^2 \in BR(\alpha^1)} u^1(\alpha)$$

game is *non-degenerate* if there is no undominated pure action  $a^2$  such that for some  $\alpha^2 \neq a^2$

$$u^i(\cdot, a^2) = u^i(\cdot, \alpha^2)$$

counterexample: player 2 gets zero always, player 1 gets either zero or one depending only on player 2's action

game is *identified* if for all  $\alpha^2$  that are not weakly dominated  
 $\rho(\cdot|\alpha^1, \alpha^2) = \rho(\cdot|\tilde{\alpha}^1, \alpha^2)$  implies  $\alpha^1 = \tilde{\alpha}^1$

$$\rho(\cdot|\alpha^1, \alpha^2) = \alpha^1 R(\alpha^2)$$

condition for identification  $R(\alpha^2)$  has full row rank for all  $\alpha^2$

## ***Patient Short Run Players: Schmidt***

short run preferences  $\begin{bmatrix} 10 & 0 \\ 0 & 1 \end{bmatrix}$

long run preferences

$$\mu^0 = 0.1 \quad \begin{bmatrix} 10 & 0 \\ 0 & 1 \end{bmatrix} \quad \text{pure coordination}$$

$$\mu^* = 0.01 \quad \begin{bmatrix} 10 & 10 \\ 0 & 1 \end{bmatrix} \quad \text{commitment type}$$

$$\mu^i = 0.89 \quad \begin{bmatrix} 1 & 1 \\ 1 & 1 \end{bmatrix} \quad \text{indifferent type}$$



strategies:

normal: play U except if you previously did D, then switch to D

commitment: always play U

indifferent type: U until deviation then D

SR: play L then alternate between R and L (on path)

if 1 deviated to D switch to R forever

if 2 deviated play L; if 1 reacts with U continue with L

reacts with D continue with R

$\delta_1 \geq .15, \delta_2 \geq .75$  then this is a subgame perfect equilibrium

- interesting deviation for SR when supposed to do R deviate to L; but then indifferent type switches to D forever
- for the normal type to prove he's not type "i" he must play D revealing he is not the commitment type

Suppose that LR can minmax SR in a pure strategy  $\underline{a}^1$

Theorem: LR gets at least  $\min_{\alpha^2 \in BR^2(\underline{a}^1)} u^1(\underline{a}^1, \alpha^2)$

let  $\underline{u}^2$  be SR minmax

let  $\tilde{u}^2$  be second best against  $\underline{a}^1$

$$N = \frac{\ln(1 - \delta_2) + \ln(\underline{u}^2 - \tilde{u}^2) - \ln(\bar{u}^2 - \tilde{u}^2)}{\ln \delta_2}$$

$$\varepsilon = \frac{(1 - \delta_2)^2 (\underline{u}^2 - \tilde{u}^2)}{(\bar{u}^2 - \tilde{u}^2)} - \delta_2^N (1 - \delta_2)$$

commit to  $\underline{a}^1$

**Lemma:** suppose  $a_2^{t+1} \notin BR(\underline{a}^1)$  with positive probability, then SR must believe that in  $t+1, \dots, t+N$  there is a probability of at least  $\varepsilon$  of not having  $\underline{a}^1$

- why is this sufficient?

*Proof of Lemma:*

2 can get at least  $\underline{u}^2$  so

$$(1 - \delta_2)u(\alpha^1, a_2^{t+1}) + \delta_2 V \geq \underline{u}^2$$

if  $pr(\underline{a}^1) > 1 - \varepsilon$  in  $t+1, \dots, t+N$

suppose the optimum is not a br; consider deviating to a br

gain at least  $\underline{u}^2 - \tilde{u}^2$  at  $t+1$ ; starting at  $t+2$  earn no less than  $\underline{u}^2$

if you hadn't deviated, starting at  $t+2$  would have earned at most

$$(1 - \delta_2) \sum_{t=1}^N \delta_2^t (\underline{u}^2 (1 - \varepsilon) + \varepsilon \bar{u}^2) + \delta_2^{N+1} \bar{u}^2$$

but we chose  $N$  and  $\varepsilon$  so that the loss exceeds the gain